

DYALOG

Glasgow 2024

Namespaces in Dyalog

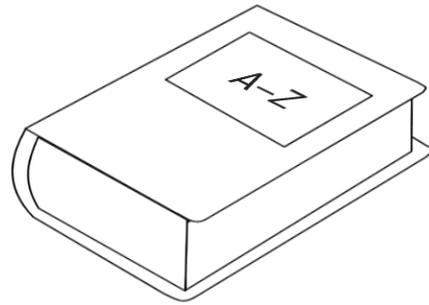


Adám Brudzewsky

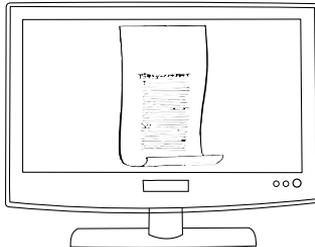
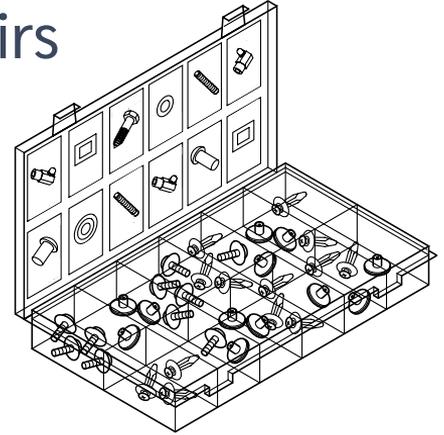


John Daintree

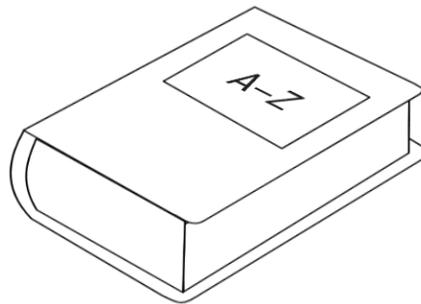
Overview



1. Self-contained dictionary of name-value pairs
2. Structured container to organise your code
3. Display form and scripted namespaces



The namespace as a
self-contained dictionary
of name-value pairs



Dictionary

```
ns ← □ NSθ  
≡ ns
```

create a namespace
it's a kind of simple scalar

0

```
]display ns (2 2 pns)
```

used like any other scalar



Dictionary

Value accessed with dot:

```
dictionary.name ← value    set value  
value ← dictionary.name    get value
```

```
ns←⎕NS⍅                    create a namespace  
ns.foo←1 2 3                populate it  
10×ns.foo                  use the value  
10 20 30
```

Dictionary: Dotted expressions

Left side of the dot must be an array of namespace(s).

```
(ns1 ns2)←(⊞NS⊖)(⊞NS⊖)  
ns1.num←77  
ns2.num←99  
ns1.num ns2.num
```

```
77 99
```

Dictionary: Dotted expressions

Left side of the dot must be an array of namespace(s):

```
(ns1 ns2).num
77 99
(2 3pns1 ns2).num
77 99 77
99 77 99
```

Dictionary: Dotted expressions

Scalar extension:

```
(a b)←55
(a b)
55 55
(ns1 ns2).num←55
(ns1 ns2).num
55 55
```

Dictionary: Dotted expressions

Scalar extension:

```
(a b)←55
(a b)
55 55
(ns1 ns2).num←55
(ns1 ns2).num ≡ (ns1.num ns2.num)
55 55
```

Dictionary: Dotted expressions

Right side of the dot can be any expression (that is, without \diamond) executed in the namespace(s) on the left side.

```
(ns1 ns2) ← (⊠NS⊠) (⊠NS⊠)
ns1.(a b) ← 'Hello' 'World'
ns1.(a,b)
HelloWorld
```

Dictionary: Dotted expressions

Right side of the dot can be any expression (that is, without \diamond) executed in the namespace(s) on the left side.

```
ns2.a←{2/ω}
ns2.b←2 2ρ6 7 8 9
ns2.(a,b)
6 6 7 7 8 8 9 9
```

Dictionary: Dotted expressions

Right side of the dot can be any expression (that is, without \diamond) executed in the namespace(s) on the left side.

`(2 3pns1 ns2).(a,b)`

<code>HelloWorld</code>	<code>6 6 7 7 8 8 9 9</code>	<code>HelloWorld</code>
<code>6 6 7 7 8 8 9 9</code>	<code>HelloWorld</code>	<code>6 6 7 7 8 8 9</code>

Exercise 1a: Creation Easy

Create a namespace req where status is 200 and Method is $\{4+2\times\omega\}$.

your \diamond expressions

req.status

200

req.Method 200

404

Exercise 1b: Updating ^{Easy}

Within `req`, apply `Method` to `status`, and update `status` with the result.

```
200      req.status
        your_expression
404      req.status
```

Exercise 1c: Importing a workspace Medium

Write a function `Into` that copies a workspace into a namespace (using `⊞CY`).

```
dfns←⊞NS⊜  
'dfns.dws' Into dfns  
dfns.disp dfns.morse 'SOS'
```



“By-value” versus “by-reference”

Non-namespaces have value semantics:

```
    A ← 1 2 3
    B ← A
    A B
1 2 3 1 2 3
```

arrays have the same value

```
    A ← 4 5 6
    A B
4 5 6 1 2 3
```

B not affected

“By-value” versus “by-reference”

Non-namespaces have value semantics:

```
A ← 1 2 3
B ← A
A B
1 2 3 1 2 3
```

arrays have the same value

```
A[3] ← 9
A B
1 2 9 1 2 3
```

B not affected

“By-value” versus “by-reference”

Non-namespaces have value semantics:

$A \leftarrow 3$

“By-value” versus “by-reference”

Non-namespaces have value semantics:

```
A ← 3  ◊  f ← A ◦ +  
f  100
```

103

```
A ← 4
```

```
f  100
```

A was passed by value – f hasn't changed

103

```
f ← A ◦ +
```

```
f  100
```

f has changed – not A

“By-value” versus “by-reference”

Non-namespaces have value semantics:

```
▽ {A}←foo A arguments are implicitly local
  A[2]←9
```

```
▽
```

```
A←1 2 3
```

```
foo A
```

```
A
```

A hasn't changed

```
1 2 3
```

“By-value” versus “by-reference”

Non-namespaces have value semantics:

```
▽ {A}←foo A arguments are implicitly local
  A[2]←9
```

```
▽
```

```
A←1 2 3
A←foo A reassigns A
```

```
A
```

```
1 9 3
```

“By-value” versus “by-reference”

Namespaces have reference semantics

- ◆ Arrays are pure values, irrespective of the actual instance
- ◆ Except any namespace, which is a container

In the context of Dyalog APL:

namespaces **are** called *references* or *refs*

Name class: 2 is array, 3 is function, 9 is scalar reference

“By-value” versus “by-reference”

Namespaces have reference semantics

```
ns1 ← NS0
ns2 ← ns1
ns2.vec
VALUE ERROR
```

both names designate the same ns
name `vec` undefined in `ns2`

```
ns1.vec ← 7 8 9
ns2.vec
7 8 9
```

name `vec` defined in `ns2`

“By-value” versus “by-reference”

Namespaces have reference semantics

```
ns1 ← □ NSθ
```

```
ns2 ← ns1
```

```
ns1.vec ← 7 8 9
```

both names designate the same ns

```
ns2.vec[2] ← 10
```

```
ns1.vec
```

value in ns1 has changed

```
7 10 9
```

“By-value” versus “by-reference”

Namespaces have reference semantics

```
ns1 ← □ NSθ  
ns1.vec ← 7 8 9  
▽ ns SetVec2 value  
  ns.vec[2] ← value  
▽
```

```
ns1 SetVec2 10  
ns1.vec
```

ns1 has been modified

```
7 10 9
```

“By-value” versus “by-reference”

Namespaces have reference semantics

```
ns1 ← NS 0
ns2 ← NS 1
ns1 . SetVec 7 8 9
ns2 . Vec[2] ← value
7 100 9
```

ns2 has been modified

```
ns3 ← NS 0
ns3 . SetVec 10
ns1 . Vec[2] ← 1000
VALUE 9 ERROR: Undefined name: vec
```

new, unrelated namespace

ns1 has been modified

“By-value” versus “by-reference”

Namespaces have reference semantics

```
ns1 ← □ NSθ
```

```
ns1.vec ← 7 8 9
```

```
▽ SetVec3 vec takes array
```

```
vec[3] ← 10
```

```
▽
```

```
SetVec3 ns1.vec
```

```
ns1.vec
```

an array: passed by value

“By-value” versus “by-reference”

How to achieve “by-value” semantics with namespaces:

```
ns1←□NSθ  
ns1.vec←7 8 9
```

```
ns3←□NS ns1  
ns3.vec[2]←10
```

clone: make “deep” copy

```
]disp (ns1 ns3).vec
```

only new copy changed



“By-value” versus “by-reference”

Namespaces have reference semantics

Allows you to:

- ◆ Track individual entities independently of their content
- ◆ Pass modifiable arguments to functions (use with care!)

“By-value” versus “by-reference”

Namespaces have reference semantics

Note the name class (40□ATX):

- depth-0 refs: 9
- other arrays: 2 (even if consisting entirely of refs)

“By-value” versus “by-reference”

Namespaces have reference semantics

Naming convention tip:

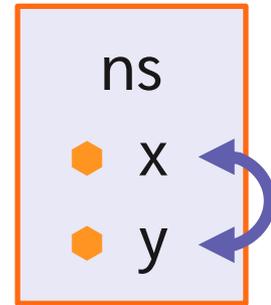
- Due to the different semantics, it may be worth naming “dottable” names in a recognisable way, for example:
`name_`

Exercise 2a: Swap names Easy

Write an *expression* that swaps the values of variables `x` and `y` in a namespace `ns`.

```
ns←⊖NS⊖  
ns.x←10 ⋄ ns.y←20  
your_expression  
ns.x ns.y
```

20 10



Exercise 2b: Swap namespaces Easy

Write an *expression* that swaps the values of the variables named `x` in the namespaces `ns1` and `ns2`.

```
ns1 ← []NSθ  ♦  ns2 ← []NSθ
```

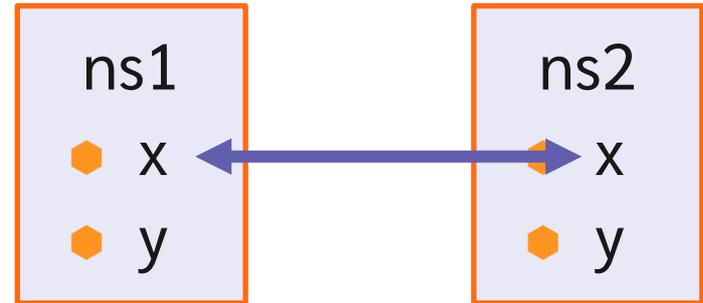
```
ns1.x ← 10  ♦  ns2.x ← 20
```

```
ns1.y ← 30  ♦  ns2.y ← 40
```

```
your_expression
```

```
ns1.x  ns1.y  ns2.x  ns2.y
```

```
20  30  10  40
```



Exercise 2c: Testing if reference Medium

Write a function `ScalarRef` that returns a scalar Boolean value indicating whether its argument is a scalar namespace.

```
ns←⊞NS⊥ ◊ ns.a←10
vec←ns.a 'abc' (ns ns) ns (⊞NS⊥) 42 ⊥ #
ScalarRef" vec
0 0 0 1 1 0 0 1
```

Use one or more of these scalar namespace properties:

- ◆ `40⊞ATX` is 9
- ◆ `⊞DR` is 326
- ◆ `Depth (≡)` is 0 (simple scalar)
- ◆ Allows dot syntax (`ns.name`)

Dictionary access under program control

```
value ← dictionary.name    get value
```

But how do we do this when the name can vary?

```
ns ← NS0  
ns.foo ← 1 2 3
```

```
name ← 'foo'  
10 × ns[name]
```

```
10 20 30
```

Dictionary access under program control

```
dictionary.name ← value    set value
```

But how do we do this when the name can vary?

```
ns←[]NSθ  
name←'foo'
```

```
ns⊕name, '←1 2 3'  
10×ns.foo
```

```
10 20 30
```

Exercise 3a: Get multiple values Medium

Write a function `Get` that takes a namespace as left argument and a vector of names as right argument and returns the corresponding values from the namespace.

```
ns←⍋NS⍅
ns.(foo bar baz)←(1 2 3) 'Hello' 42

names←'foo' 'bar' 'foo' 'baz'
ns Get names
1 2 3 Hello 1 2 3 42
```

Exercise 3b: Set a value Medium

Write a function `Set` that takes a two-element `(name value)` vector as right argument, then does the corresponding assignment.

```
      Set 'my' (15)
      10×my
10 20 30 40 50
```

Exercise 3c: Set in a namespace Medium

Improve `Set` so that it takes a namespace as left argument and does the corresponding assignment in that namespace.

```
ns ← NS0
ns Set 'my' (15)
10×ns.my
10 20 30 40 50
```

Exercise 3d: Set multiple values ^{Hard}

Improve `Set` so that it handles multiple two-element `(name value)` vectors as right argument.

```
ns ← □ NSθ
ns Set ('yours' 'Hello') ('mine' 'World')
ns.(yours mine)
Hello World
```

Dictionary access under program control

Beware: \oplus is potentially dangerous.

```
ns ← □NSθ  
name ← ' □OFF◇ '      !!!  
value ← 1 2 3  
  
ns {α. {⊕α, '←ω'} / ω} name value
```

Solution: In production, validate your arguments ($\neg 1 \neq 40 \square ATX$).
— Look forward to $\square VGET$ and $\square VSET$ in version 20.0...

Dictionary access in version 20.0

Exercise 3a: Get multiple values

`Get ← □VGET`

Exercise 3b: Set a value

`Set ← {□VSET c ω} or Get ← □VSET c`

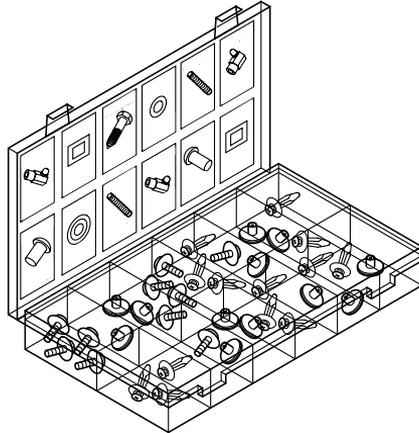
Exercise 3c: Set in a namespace

`Set ← {α□VSET c ω} or Set ← □VSET o c`

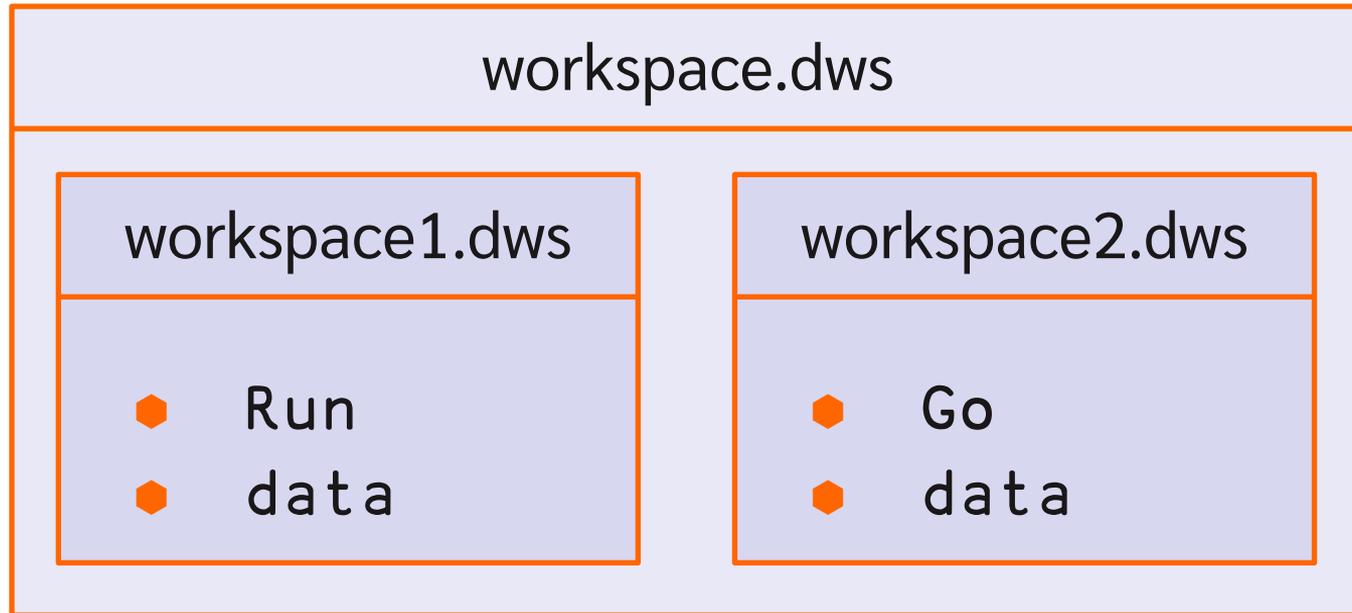
Exercise 3d: Set multiple values

`Set ← □VSET`

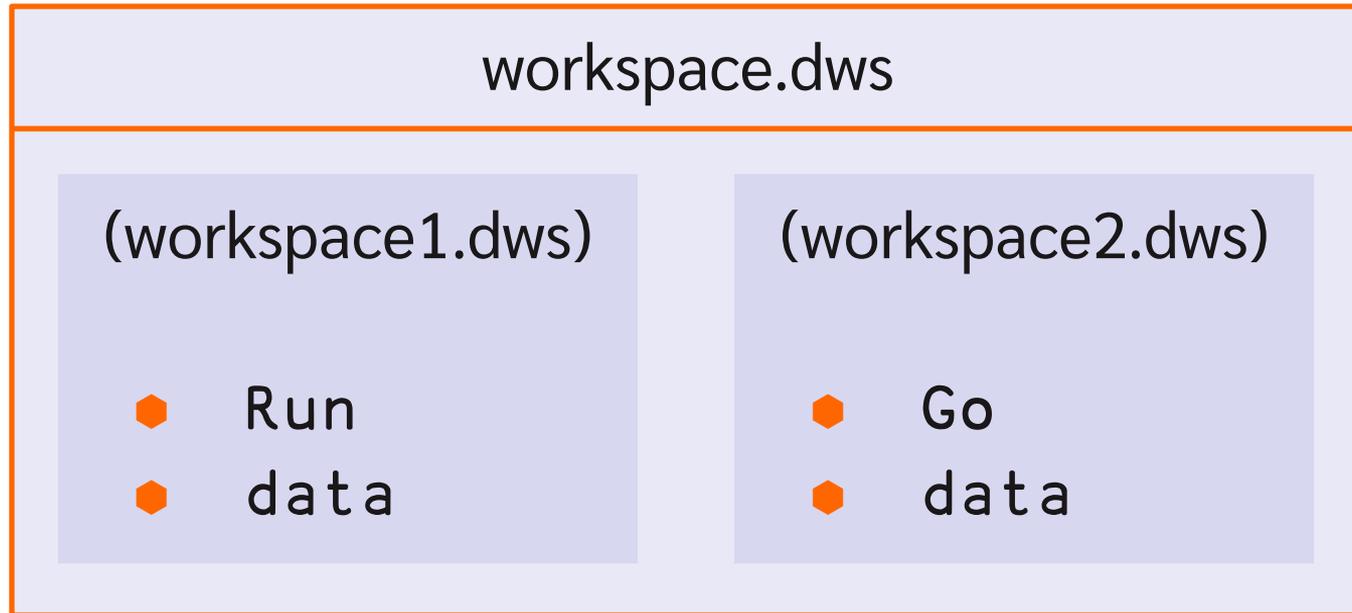
The namespace as a
structured container
to organise your code



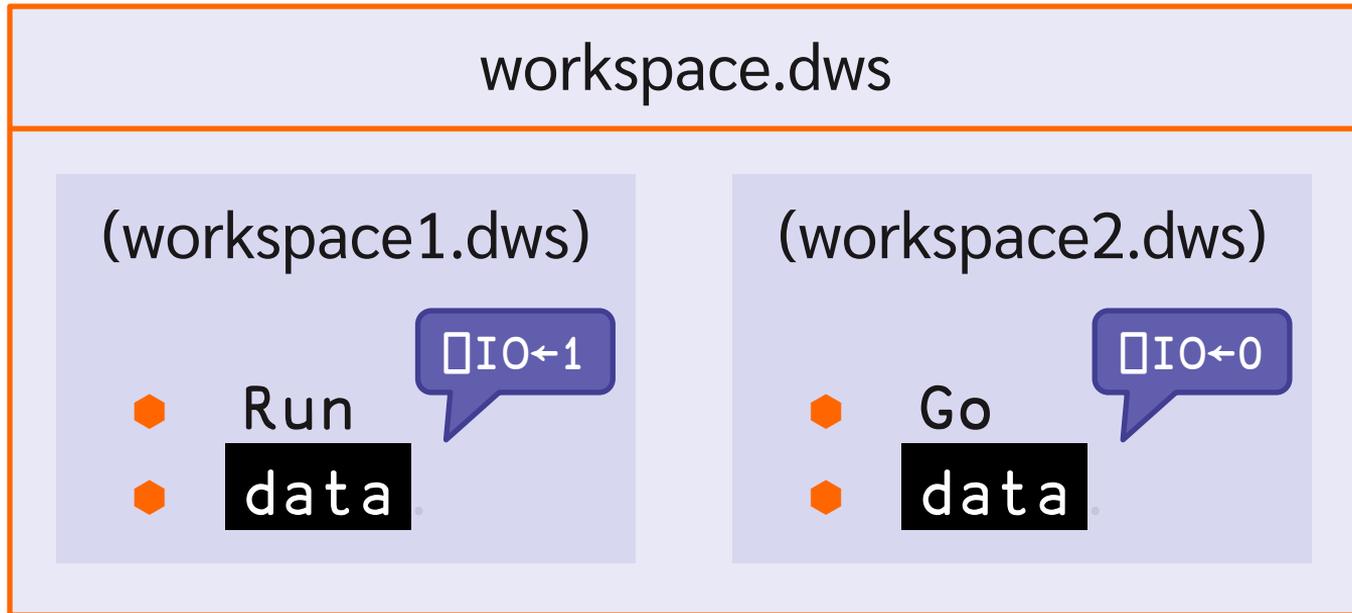
Structured container



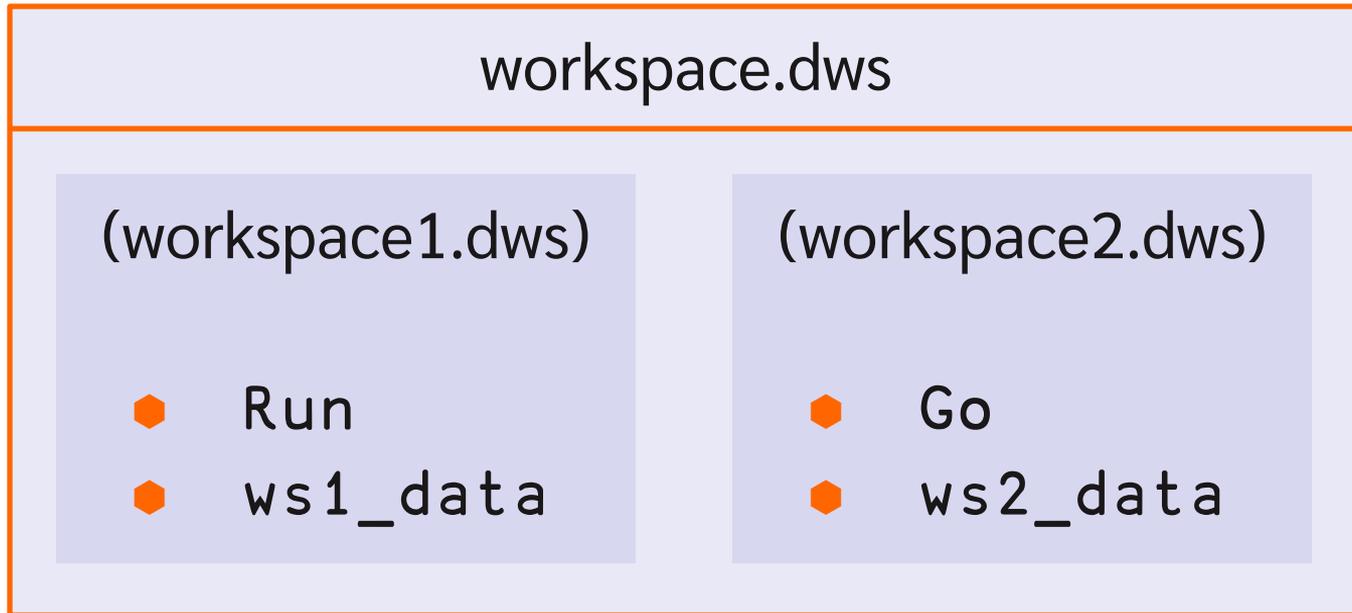
Structured container



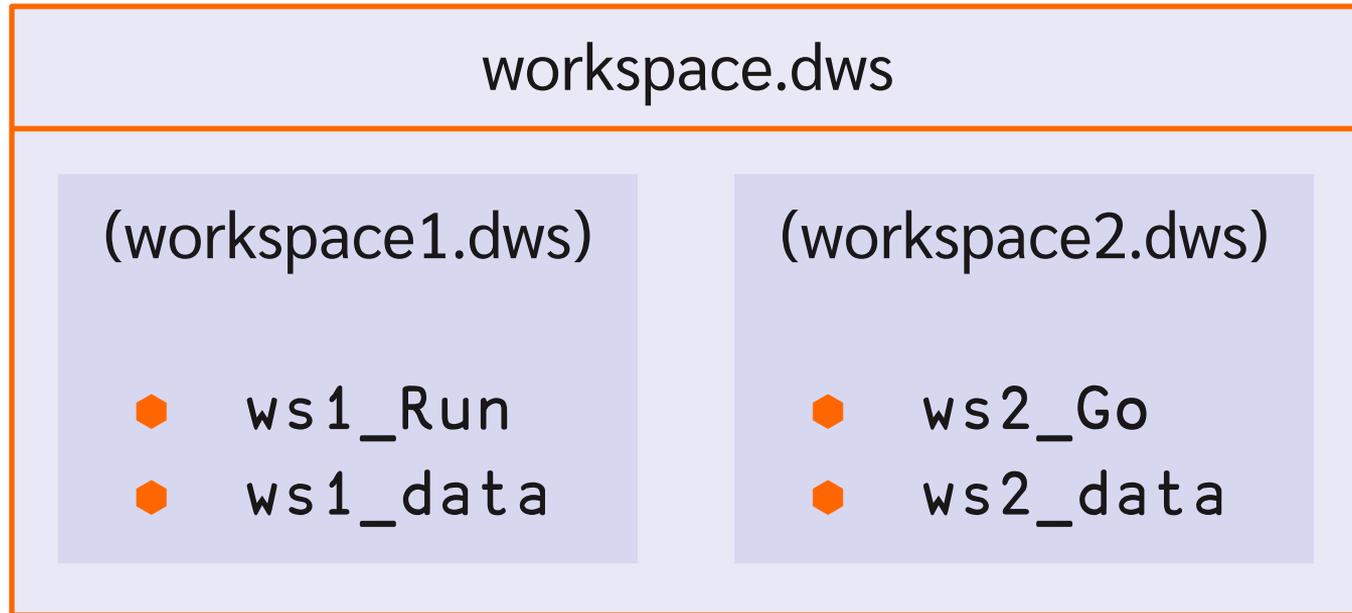
Structured container



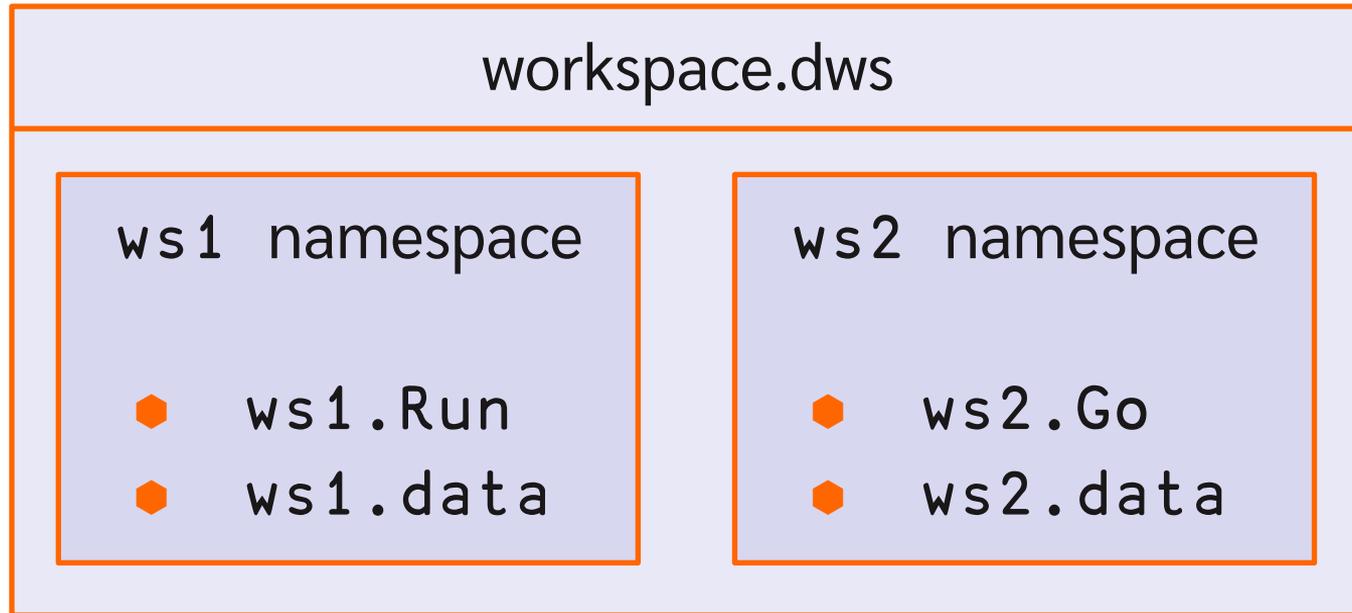
Structured container



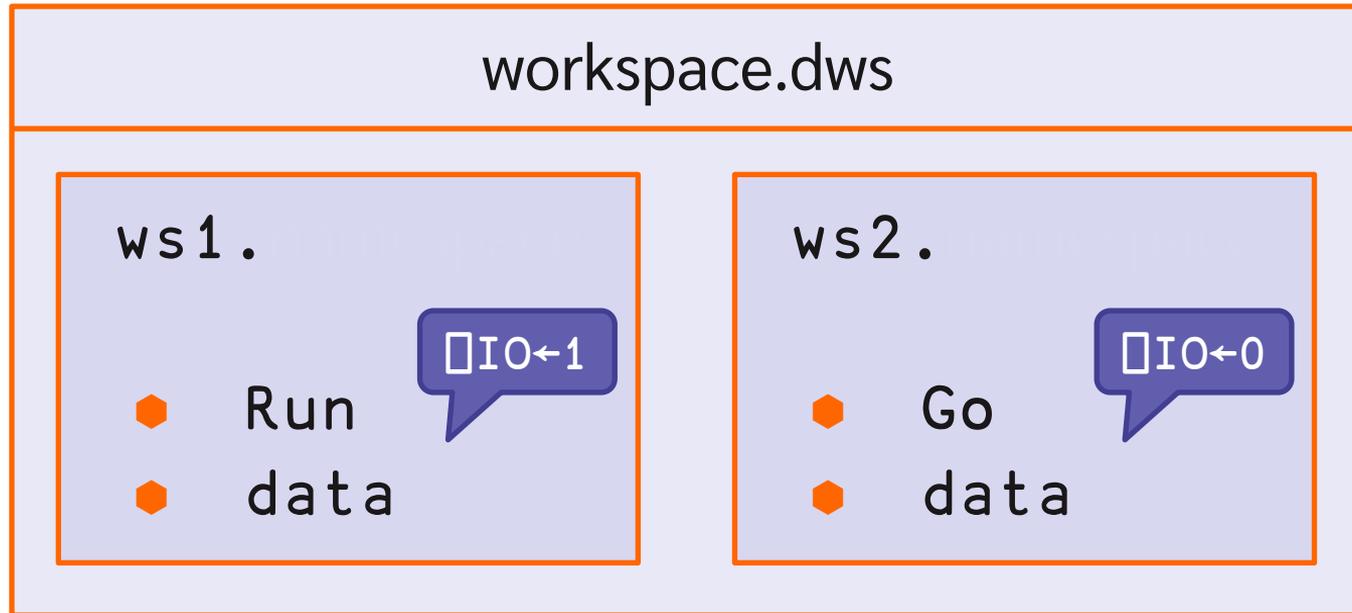
Structured container



Structured container



Structured container



Structured container

Tree

Node

Leaf

Roots

Separator

Parent node

Access current node

Report current node

Workspace

Namespace

Name

□SE

.

##

□THIS

□THIS

File system

Directory

File

UNIX

/

/

..

.

pwd

WINDOWS

C:\ D:\

\

..

.

cd

Structured container: actions

Tree

Node

Leaf

Create node

Copy node

Change current node

List node contents

⋮

Workspace

Namespace

Name

□NS

□NS

□CS

□NL

⋮

File system

Directory

File

UNIX

mkdir

cp

cd

ls

⋮

WINDOWS

md

copy

cd

dir

⋮

Structured container

```
data←7 8 9
ns←⊞NS⊜
ns.data←'hello'
(data)(ns.data)
```



```
⊞CS ns      Change Space
⊞THIS
#. [Namespace]
  data
hello
⊞CS ##      Return to parent
⊞THIS
#
  data
7 8 9
```

Structured container

```
data←7 8 9
ns←⊞NS⊜
ns.data←'hello'
(data)(ns.data)
```



```
ns.(sub←⊞NS⊜)
ns.sub.data←'x'
```

```
⊞CS ns      Change Space
data
hello
##.data
7 8 9
⊞CS sub    Change Space
data
x
##.data
hello
```

Structured container

```
data←7 8 9
ns←⊞NS⊜
ns.data←'hello'
(data)(ns.data)
```

7 8 9	hello
-------	-------

```
ns.(sub←⊞NS⊜)
ns.sub.data←'x'
```

```
⊞CS ns Change Space
sub.(data #.data)
```

x	7 8 9
---	-------

Exercise 4a: Is argument a root? Easy

Write a function `IsRoot` that takes a namespace as argument that tells you whether that namespace is a root namespace.

```
0      IsRoot []SE.Dyalog.Util s
1      IsRoot #
1      IsRoot []SE
```

Exercise 4b: What is my root? Medium

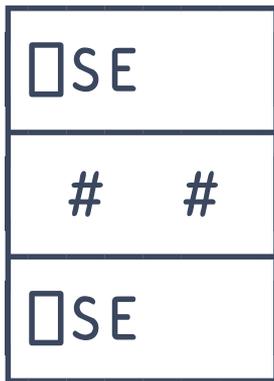
Write a function `FindRoot` that takes a namespace as argument and returns its root.

```
FindRoot []SE.Dyalog.Utils
[]SE
FindRoot #
#
FindRoot []NS0
#
```

Exercise 4c: Our roots? Easy — based on FindRoot

Write a function `FindRoots` that takes an arbitrary array of namespaces and finds the root for each namespace.

```
FindRoots ;[SE.Dyalog.Utils(#,[NSθ])SE
```



Exercise 4d: Namespace lineage ^{Hard}

Write a function `Line` that takes a single namespace and returns its lineage (as a vector of refs) from root to leaf.

```
Line []SE.Dyalog.Utils
[]SE []SE.Dyalog []SE.Dyalog.Utils

Line []SE.cbbot.bandsb2.sb.io
[]SE []SE.cbbot []SE.cbbot.bandsb2
[]SE.cbbot.bandsb2.sb []SE.cbbot.bandsb2.sb.io
```

Parent hierarchy

Namespaces have a single immutable parent (fixed at creation time) which is the namespace where `⊠NS` was called.

```
ns←⊠NS⊘
```

```
#=ns.##
```

created in #

1

```
ns.sub1←ns.⊠NS⊘
```

```
ns=ns.sub1.##
```

created in #.ns

1

Parent hierarchy

Namespaces have a single immutable parent (fixed at creation time) which is the namespace where `NS` was called.

```
ns.sub1.(sub2←NSθ)
ns.sub1=ns.sub1.sub2.##    created in #.ns.sub1
1
ns.sub1.sub2.##.##.##
#
```

Parent hierarchy

```
ns0 ← □ NSθ  
ns1 ← □ NSθ  
ns1 . ns2 ← ns0
```

Names in a namespace are **not** (necessarily) its children:

```
ns1 = ns1 . ns2 . ##
```

0

Note: `ns1 . ns2` \equiv `ns0` which was created in `#`

Parent hierarchy

```
ns0 ← □ NSθ  
ns1 ← □ NSθ  
ns1.ns2 ← ns0
```

Children are **not** (necessarily) named from their parent:

```
ns3 ← ns1.ns2.□ NSθ
```

Note: there exists no name such that $ns3 \equiv ns3.##\underline{\text{name}}$:

```
ns1.ns2.ns3
```

```
VALUE ERROR: Undefined name: ns3
```

Two approaches to namespaces:

Purpose:	Code structure	Data dictionary
Usage:	rooted	self-contained
Creation:	named	unnamed
Display:	full path	description
Hierarchy:	tree	flat

Each approach is easy, mixing them is tough — don't try!