

ALG ϕ RITHMS AS A T $\phi\phi$ L ϕ F TH ϕ UGHT

Conor Hoekstra



code_report





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code_report



Algorithms +
Data
Structures =
Programs



RAPIDS





RAPIDS





1.



2.



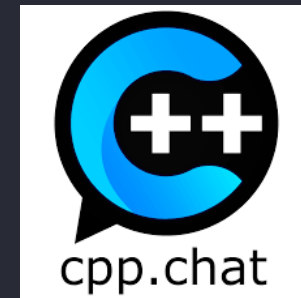
3.



4.



5.





1979 ACM Turing Award Lecture

Delivered at ACM '79, Detroit, Oct. 29, 1979

The 1979 ACM Turing Award was presented to Kenneth E. Iverson by Walter Carlson, Chairman of the Awards Committee, at the ACM Annual Conference in Detroit, Michigan, October 29, 1979.

In making its selection, the General Technical Achievement Award Committee cited Iverson for his pioneering effort in programming languages and mathematical notation resulting in what the computing field now knows as APL. Iverson's contributions to the implementation of interactive systems, to the educational uses of APL, and to programming language theory and practice were also noted.

Born and raised in Canada, Iverson received his doctorate in 1954 from Harvard University. There he served as Assistant Professor of Applied Mathematics from 1955-1960. He then joined International Business Machines, Corp. and in 1970 was named an IBM Fellow in honor of his contribution to the development of APL.

Dr. Iverson is presently with I.P. Sharp Associates in Toronto. He has published numerous articles on programming languages and has written four books about programming and mathematics: *A Programming Language* (1962), *Elementary Functions* (1966), *Algebra: An Algorithmic Treatment* (1972), and *Elementary Analysis* (1976).

Notation as a Tool of Thought

Kenneth E. Iverson
IBM Thomas J. Watson Research Center



Key Words and Phrases: APL, mathematical notation
CR Category: 4.2

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The importance of nomenclature, notation, and language as tools of thought has long been recognized. In chemistry and in botany, for example, the establishment of systems of nomenclature by Lavoisier and Linnaeus did much to stimulate and to channel later investigation. Concerning language, George Boole in his *Laws of Thought* [1, p.24] asserted "That language is an instrument of human reason, and not merely a medium for the expression of thought, is a truth generally admitted."

Mathematical notation provides perhaps the best-known and best-developed example of language used consciously as a tool of thought. Recognition of the important role of notation in mathematics is clear from the quotations from mathematicians given in Cajori's *A History of Mathematical Notations* [2, pp.332,331]. They are well worth reading in full, but the following excerpts suggest the tone:

By relieving the brain of all unnecessary work, a good notation sets it free to concentrate on more advanced problems, and in effect increases the mental power of the race.

A.N. Whitehead

ϕ \emptyset \cup Δ \subseteq \uparrow \wr \in

ϕ \emptyset \cup ∇ \subseteq \uparrow \wr \in
•• / \ ~ \boxplus \ast • \ddot{o}

ALGORITHMS

ϕ \emptyset \cup ∇ \subseteq \uparrow \wr \in
.. / \ \approx \boxplus \ast \cdot $\ddot{\circ}$

OPERATORS

FUNCTIONS

ϕ \emptyset \cup \cap \subseteq \uparrow \wr \in
.. / \ \approx \equiv \times \cdot \circ

OPERATORS

VERBS

ϕ \emptyset \cup ∇ \subseteq \uparrow \wr \in
¨ / \ ~ \boxplus \ast \cdot $\ddot{\circ}$

ADVERBS & CONJUNCTIONS

ALGORITHMS

ϕ \emptyset \cup \cap \subseteq \uparrow \downarrow \in
.. / \ \approx \equiv \times \cdot \circ

OPERATORS

ALGORITHMS

ϕ reverse

υ unique

ξ partition

ι iota

ALGORITHMS



ϕ	reverse	reverse
u	unique	unique
\subseteq	partition	partition
i	iota	iota

allEqual



```
import Data.List.HT (allEqual)
```

```
allEqual [1,2,3,4] -- False
```

```
allEqual [1,1,1,1] -- True
```



```
from more_itertools import all_equal
```

```
all_equal([1,2,3,4]) # False
```

```
all_equal([1,1,1,1]) # True
```

Hoogle Translate

allEqual



Python

all_equal

more-itertools

[Doc](#)



JavaScript

allEqual

bbo

[Doc](#)



Haskell

allEqual

Data.List.HT

[Doc](#)



Clojure

apply =

[Doc](#)



Racket

apply =

[Doc](#)

SOLUTION #1



^ / 3 = 1



first



⊃ 1 2 3 4



1



$\{ (\supset \omega) = \omega \}$ 1 2 3 4



{ (=> 1 2 3 4) = 1 2 3 4 }



1 = 1 2 3 4



1 0 0 0



















^ / 1 0 0 0



^ / 1 0 0 0

1 ^ 0 ^ 0 ^ 0

	APL	/ (reduce)	-	Doc
	CUDA	reduce	Thrust	Doc
	D	reduce	algorithm.iteration	Doc
	Ruby	reduce	Enumerable	Doc
	Python	reduce	itertools	Doc
	Elixir	reduce	Enum	Doc
	Kotlin	reduce	collections	Doc
	Clojure	reduce	core	Doc
	C++	reduce	<numeric>	Doc
	Haskell	foldl	Data.List	Doc
	Racket	foldl	base	Doc
	Rust	fold	trait.Iterator	Doc
	q	over	-	Doc
	C#	Aggregate	Enumerable	Doc
	J	/ (insert)	-	Doc
	C++	accumulate	<numeric>	Doc



^ / 1 0 0 0

1 ^ 0 ^ 0 ^ 0



0



$$\wedge / \{ (\sup \omega) = \omega \}$$



$$\wedge / \{ (\supset \omega) = (\vdash \omega) \}$$



^ / 3 = 1



^ / ≡ ⊢

⊢ first

= equal

^ and

⊢ same

/ reduce

SOLUTION #2



$(1 = \neq) \cup$



U

unique



u 1 2 2 4



1 2 4



≠ 1 2 4

tally



3



$$1 = 3$$



0



1 = ~~≠~~ ∪ 1 2 2 4



FREE

$(1 = \neq) \cup$



(1 = ≠) u

u unique
= equal
≠ tally



(1 ≥ ≠) ∪

∪ unique

≥ gte

≠ tally

SOLUTION #3



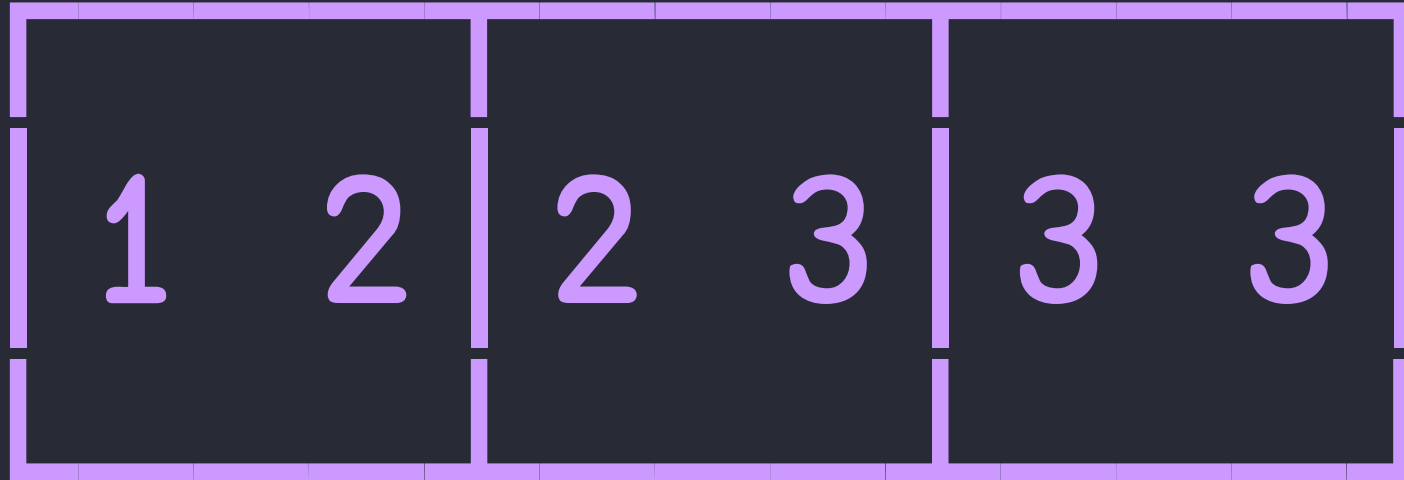
$$\wedge / 2 = / \vdash$$



2 = / 1 2 3 3

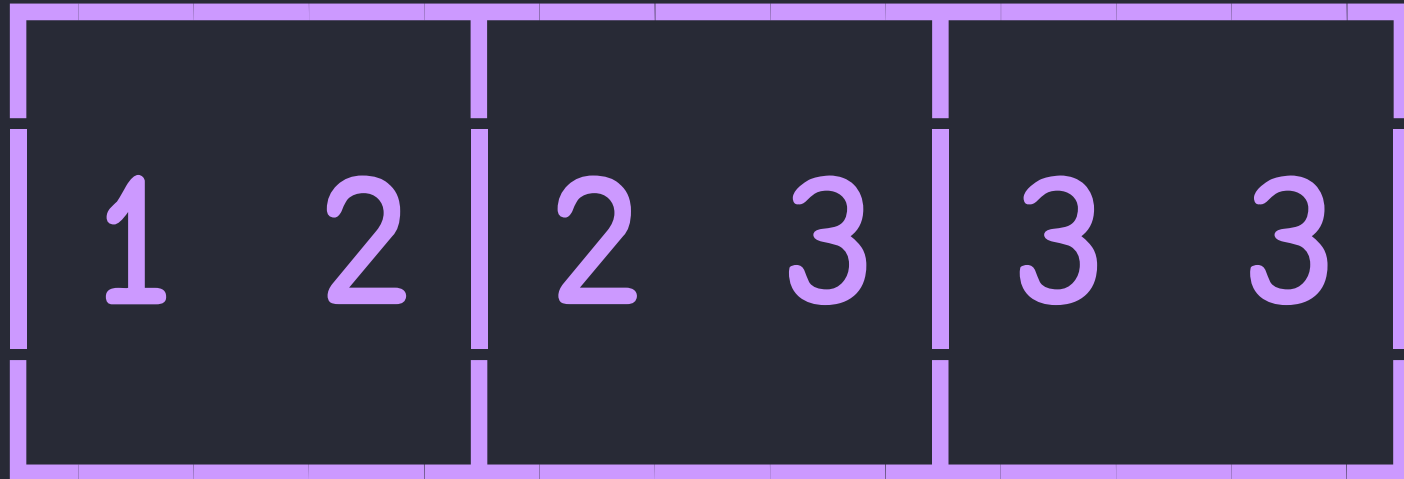


2, / 1 2 3 3



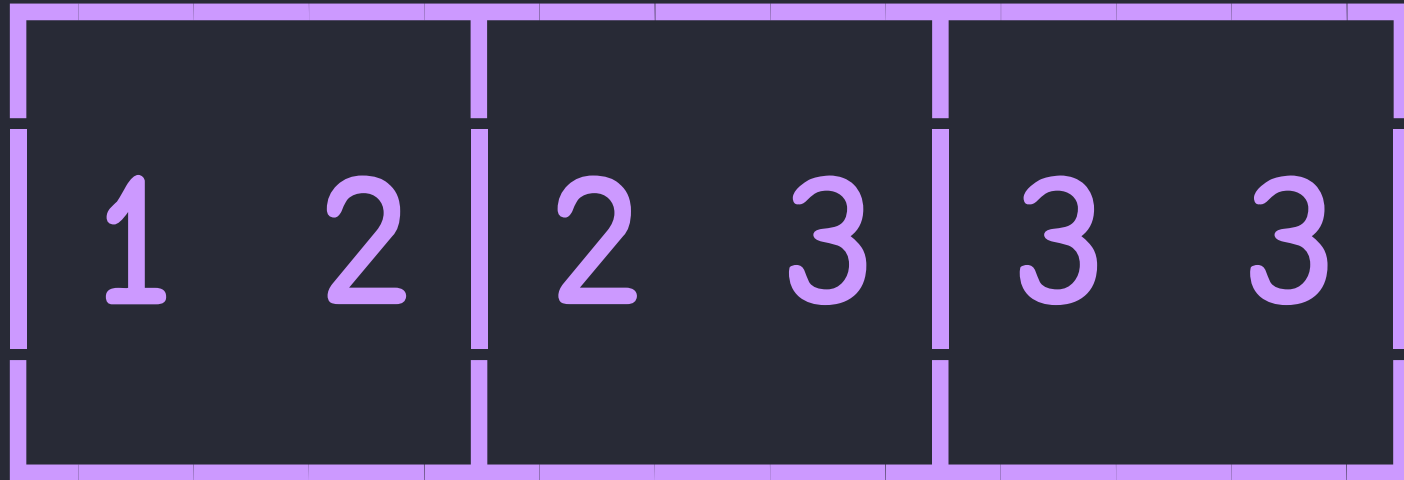


2, / 1 2 3 3





2 = / 1 2 3 3





2 = / 1 2 3 3

1 = 2	2 = 3	3 = 3
-------	-------	-------



2 = / 1 2 3 3

0	0	1
---	---	---



2 = / 1 2 3 3

1 = 2	2 = 3	3 = 3
-------	-------	-------



2 + / 1 2 3 3


1 + 2	2 + 3	3 + 3
-------	-------	-------



2 + / 1 2 3 3

3	5	6
---	---	---



`pandas.DataFrame.rolling()` 
`cudf.Series.rolling()` 



`cudf::rolling_window()` 



2 = / 1 2 3 3



0 0 1



^ / 0 0 1



0



$$\wedge / 2 = / \vdash$$



$\wedge / 2 = / \vdash$

\wedge and

$=$ equal

\vdash same

$/$ reduce

SOLUTION #4



L / = r /



L / 1 2 3 4



1



L / 1 2 3 4



┌ / 1 2 3 4



4



$$\{ (L / \omega) = \lceil \omega \}$$



$$\{ (L / \omega) = (\Gamma / \omega) \}$$



L / = r /



L / = ⌈ /

L min

⌈ max

= equal

/ reduce



/	reduce	3
=	equal	3
^	and	2
⊢	same	2
⊃	first	1
∪	unique	1
≥	gte	1
≠	tally	1
⊥	min	1
⌈	max	1

SOLUTION #5



(1 ≥ ≠) → 目



{ $\alpha\omega$ }  7 8 8 9 9 9



{ $\alpha\omega$ } 7 8 8 9 9 9

7	1		
8	2	3	
9	4	5	6



{ α } \exists 7 8 8 9 9 9

7 8 9



+ 7 8 8 9 9 9

7 8 9



(1 ≥ ≠) → 7 8 8 9 9 9

0



(1 ≥ ≠) → 目

SOLUTION #6



3 ^ . = +



(\supset \wedge \cdot $=$ \vdash) 1 2 3



$$(\supset 1 \ 2 \ 3) \wedge . = (\vdash 1 \ 2 \ 3)$$



$(\Rightarrow 1 \ 2 \ 3) \wedge . = 1 \ 2 \ 3$



1 ^ . = 1 2 3



$\wedge \cdot =$

1 1 1

1 2 3



$\wedge \cdot =$

1 0 0



$\wedge \cdot =$

0



SOLUTIONS

#1 $\wedge / \supset = \vdash$

#2 $(1 = \neq) \cup$

#3 $\wedge / 2 = / \vdash$

#4 $\lfloor / = \lceil /$

#5 $(1 \geq \neq) \dashv \equiv$

#6 $\supset \wedge \cdot = \vdash$

#7 $\uparrow \equiv \psi$

#8 $1 \circ \phi \equiv \vdash$

⌘ Aaron Hsu

⌘ Adám Brudzewsky

⌘ Bob Therriault



/	reduce	3
=	equal	3
^	and	2
⊢	same	2
⊃	first	1
∪	unique	1
≥	gte	1
≠	tally	1
⊥	min	1
⌈	max	1



=	equal	4
⊢	same	4
/	reduce	3
^	and	3
⊃	first	2
≥	gte	2
≠	tally	2
∪	unique	1
⌊	min	1
⌈	max	1
→	left	1
⊟	key	1
.	inner	1

ALGORITHMS AS A TOOL OF THOUGHT

=	equal	4
⊢	same	4
/	reduce	3
^	and	3
⊃	first	2
≥	gte	2
≠	tally	2
∪	unique	1
⊥	min	1
⌈	max	1
→	left	1
⊞	key	1
.	inner	1


```

      3 3ι9
3
      3 3ρι9
1 2 3
4 5 6
7 8 9
      {αω}⊖1 1 2 2


|   |   |   |
|---|---|---|
| 1 | 1 | 2 |
| 2 | 3 | 4 |


      ≠∘ι⊖1 1 2 2
1 1
      ι∘≠⊖1 1 2 2
2 2
      ι10
1 2 3 4 5 6 7 8 9 10
      APL

```

```

APL x
[0] APL←{
[1]   □←'APL IS AWESOME'
[2]   □←'RIDE IS AWESOME'
[3]   }

```

SI: 0 &: 1 Ln 1, Col 13

```
3 3 19
3
1 2 3
4 5 6
7 8 9
{αω}⊖1 1 2 2


|   |   |   |
|---|---|---|
| 1 | 1 | 2 |
| 2 | 3 | 4 |


≠⊖1 1 2 2
1 1
┌≠⊖1 1 2 2
2 2
ι10
1 2 3 4 5 6 7 8 9 10
APL
APL←{
  [0] 'APL IS AWESOME'
  [1] 'RIDE IS AWESOME'
  [2]
  [3] }
```

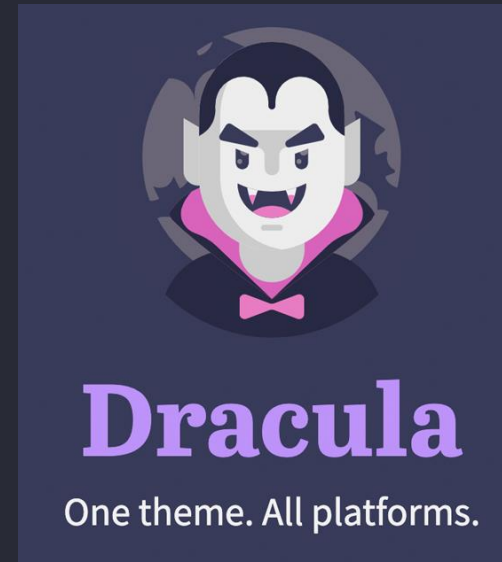
<https://www.dyalog.com/>
<https://github.com/Dyalog/ride>

```
3 3 9
3
3 3 9
1 2 3
4 5 6
7 8 9
{αω}⊖1 1 2 2


|   |   |   |
|---|---|---|
| 1 | 1 | 2 |
| 2 | 3 | 4 |


≠⊖1 1 2 2
1 1
⊖≠⊖1 1 2 2
2 2
ι10
1 2 3 4 5 6 7 8 9 10
APL
```

```
[0] APL ← {
[1]   ← 'APL IS AWESOME '
[2]   ← 'RIDE IS AWESOME '
[3] }
```



<https://www.dyalog.com/>
<https://github.com/Dyalog/ride>

Thank You!

Conor Hoekstra



code_report



NVIDIA

RAPIDS

#include

Questions?

Conor Hoekstra



code_report



nVIDIA

RAPIDS

#include